MODEL-BASED TESTING AND MODEL CHECKING FOR SAFETY-CRITICAL HIERARCHICAL SYSTEMS

December 6, 2018

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THIS THESIS WORK AN OVERVIEW



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 - a new test case generation procedure;
 - a logic to express linear properties of concurrent hierarchical computations;
 - a model checking procedure for hierarchical machines;



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- 6. A model checking procedure for hierarchical machines.

SYSTEMS VERIFICATION

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The process of checking that a system meets certain requirements derived from a given *specification*.

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Why should we care?

- Computer systems are ubiquitous and we depend more and more on them;
- Malfunctions may cause financial losses or worse!

SYSTEMS VERIFICATION WHEN CLASSIC TECHNIQUES FALL SHORT



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- ► ineffective with complex concurrent systems;
- expensive.

SYSTEMS VERIFICATION THE MODEL-BASED APPROACH



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The model-based approach integrates formal methods in the verification process.

abstract a model (formal specification) of the system;

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 - model checking.

MODEL CHECKING THE MODEL-BASED APPROACH



Model checking

Given a model and a property, check that the property holds in **every possible** model behaviour.

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G CONTRACTOR

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produces counter-examples;

• can be integrated earlier in the development process.

PREVIOUS RESEARCH

THE CRYSTAL PROJECT CRITICAL SYSTEMS ENGINEERING ACCELERATION





- driven by real-world industrial use cases from the automotive, aerospace, rail and health sector;
- produce or improve tool chains to reduce design/verification costs.



PREVIOUS RESEARCH WITHIN THE CRYSTAL PROJECT



Introduced Dynamic State Machines (DSTM);

PREVIOUS RESEARCH WITHIN THE CRYSTAL PROJECT

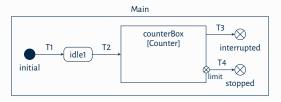


- Introduced Dynamic State Machines (DSTM);
- devised an automatable test case generation procedure for DSTM;

DYNAMIC STATE MACHINES

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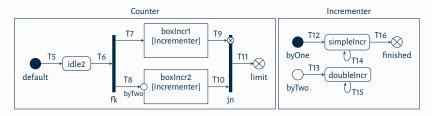
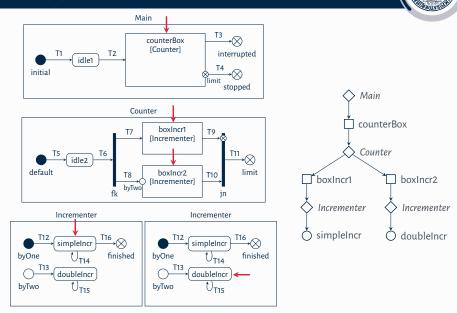


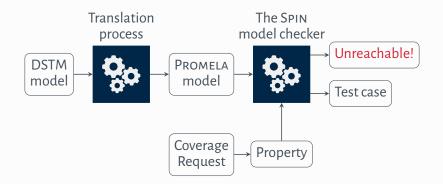
Figure: A simple DSTM specification

DYNAMIC STATE MACHINES

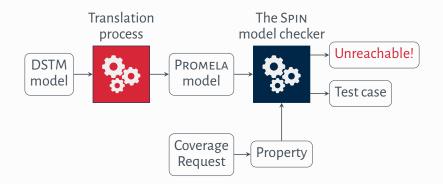


MODEL-BASED TESTING ON DSTM MODELS

MODEL-BASED TESTING GENERATING TEST-CASES FROM DSTM MODELS



MODEL-BASED TESTING GENERATING TEST-CASES FROM DSTM MODELS



HIERARCHICAL LINEAR-TIME TEMPORAL LOGIC WITH INTERRUPTS

LINEAR-TIME TEMPORAL LOGIC (LTL) Reasoning about linear time

$$\phi \coloneqq p \in \mathcal{P} \mid \neg \phi \mid \phi \lor \phi$$

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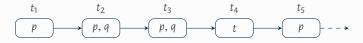
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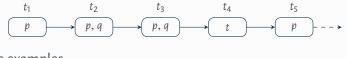


Some examples:

• $X(p \land q)$ (Next)

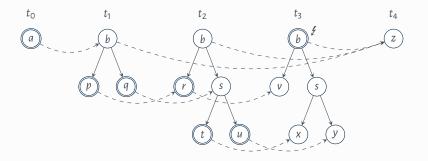
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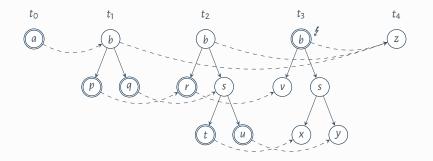


Some examples:

HIERARCHICAL LTL WITH INTERRUPTS A NOVEL EXTENSION OF THE WELL-KNOWN LTL

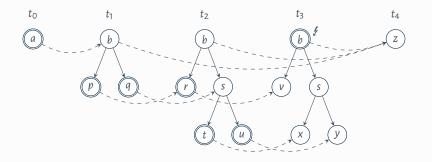


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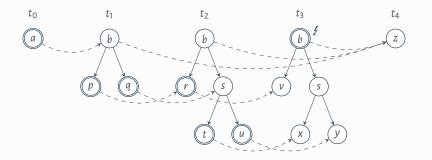
► LTL-like operators to reason about the evolution of a module;

HIERARCHICAL LTL WITH INTERRUPTS A NOVEL EXTENSION OF THE WELL-KNOWN LTL



- LTL-like operators to reason about the evolution of a module;
- new operators to contextualize formulae in the hierarchy of modules.

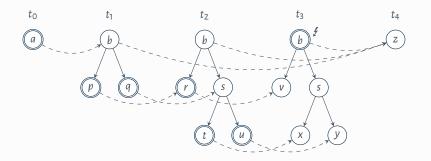
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New operators:

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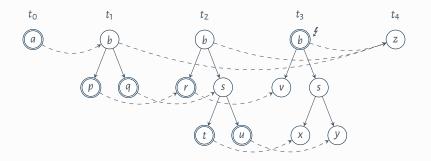
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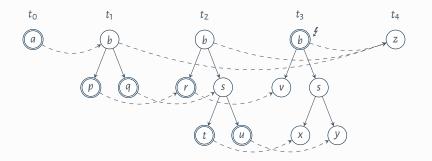
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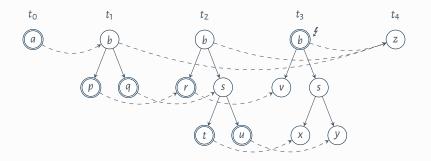


HIERARCHICAL LTL WITH INTERRUPTS A NOVEL EXTENSION OF THE WELL-KNOWN LTL





HIERARCHICAL LTL WITH INTERRUPTS A NOVEL EXTENSION OF THE WELL-KNOWN LTL



•
$$X_{f}(z);$$

►
$$X(\downarrow_1(p \land \rightarrow(q)));$$

COMMUNICATING HIERARCHICAL AUTOMATA

$\begin{array}{l} \text{SIMPLIFYING DSTM} \\ \text{Communicating Hierarchical Automata with Interrupts (CHA^{\sharp})} \end{array}$

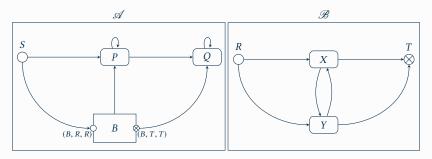


Figure: A simple CHA[∉]

SIMPLIFYING DSTM Communicating Hierarchical Automata with Interrupts (CHA⁴)

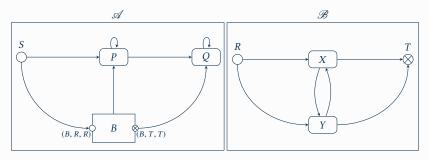


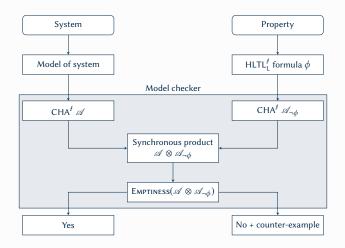
Figure: A simple CHA[£]

Differences with DSTM:

- no recursion;
- no forks/joins;

MODEL CHECKING

THE MODEL CHECKING PROCEDURE OVERVIEW



CONCLUSIONS

CONCLUSIONS Achieved goals and further research



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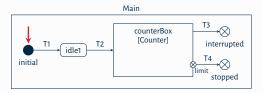
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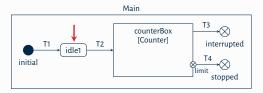
- extend results to DSTM;
- implement model checking in the existing DSTM tool chain.

Thank you for your time!

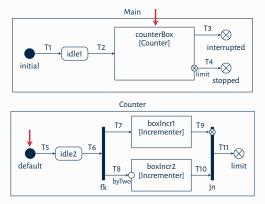
BACKUP SLIDES

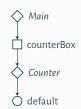


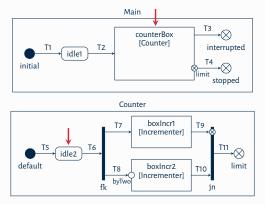


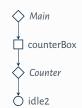


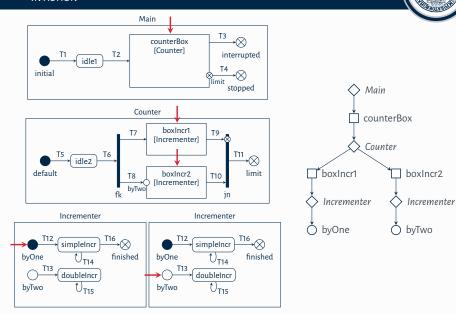


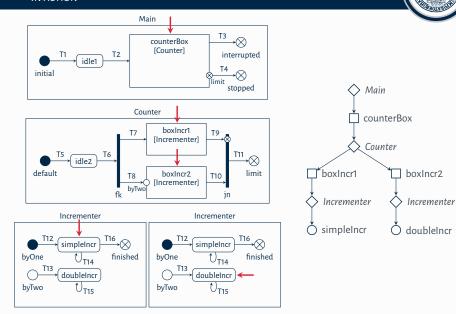


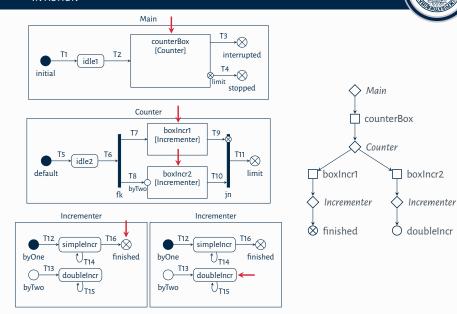


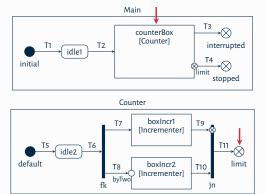


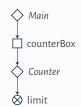


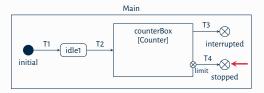








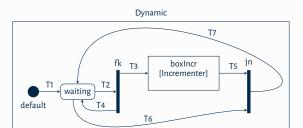






DYNAMIC STATE MACHINES A DYNAMIC INSTANTIATION EXAMPLE





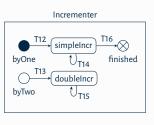


Figure: A dynamic DSTM specification

Deciding Emptiness for CSA^{\sharp}

THE EMPTINESS PROCEDURE



Emptiness problem

Given a CHA^{t} , is there any accepting computation?

Main intuition:

bottom-up summarization;

DECIDING SATISFIABILITY FOR HLTL[#] FORMULAE

THE SATISFIABILITY PROCEDURE

Satisfiability problem

Given a HLTL^f formula ϕ , is there any hierarchical computation satisfying ϕ ?

Main intuitions:

- ▶ build a CHA^f accepting hierarchical words satisfying the formula;
- non-trivial extension of the Vardi-Wolper approach for LTL;
- check for the emptiness of the automaton;

THE MODEL CHECKING PROCEDURE

